

Experiments on the Generation of Distinguishing N-Tuples for Selected Character Dichotomies

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1 Introduction

This report is a companion to [1]. In the present document, we describe experiments on generating n-tuples for optical character recognition. The focus is on the generation process itself, and not the use of the generated tuples. The experiments were conducted at Rensselaer during late January, 1995.

An outline of the report is as follows. In Section 2 we describe the experiments. The presentation of the experiment data is explained in Section 3. Section 4 contains some discussion. The experiment data is presented in Appendix A.

This report assumes familiarity with [1]. Additional information regarding the generators, and further discussion of the experiments, may be found there.

2 Description of Experiments

We consider the generation of n -tuples using two generators, called Gen0 and Gen1. Gen0 is a simple backtracking algorithm. Gen1 is a more sophisticated backtracking algorithm.

The main purpose of the experiments was to show the following.

1. Distinguishing tuples can be found reasonably quickly, when they exist.
2. Gen1 is a reasonably efficient algorithm for generating tuples. This was to be shown by comparing Gen1 to the benchmark program Gen0.
3. The execution time of Gen1 can be controlled by varying the difficulty of the problem as defined by the p and q parameters.

We also consider the quality of solution tuples as a function of the dichotomy.

The experiments consisted of executing Gen0 and Gen1 on selected character dichotomies for $n = 4, 7, 10$. Five dichotomies were used, giving $2 \cdot 3 \cdot 5 = 30$ experiments. For each experiment, Appendix A contains one table summarizing the results. Combined, the experiments represent more than 40,000 generator invocations, and hundreds of hours of CPU time. (Unless otherwise specified, all times in this document are CPU times.)

The experiments were conducted on 4 two-processor SPARC 20's, using one CPU at a time. Three of the SPARC 20's had 64 megabytes of memory, and one had 256 megabytes. The generation process had modest memory requirements, e.g., a megabyte or less, so the amount of memory possessed by the machines was not critical for these experiments.

The generators were executed in the mode "find a solution for a specified (p, q) pair." The generators executed until either a solution was found, or the search terminated because a maximum search node limit was reached. The maximum node limits were selected so that the generators executed for

about 5 minutes on a certain benchmark problem, when no solution was found. This “failure time” is different for different problems, because the time spent at a given search node varies. For these experiments, failure times were usually between 2 and 25 minutes. In some cases, when the problem was extremely constrained as when q is very small, failure times were only a few seconds.

For a given dichotomy, value of n , and generator, the generator was executed according to the following C-like pseudocode.

```
integer p,q,s,f; // for given p,q: s=successes, f=failures

for (q = n-1, p = 0, s = 50; (q >= 2) and (p != 1); q = q-1) {
  for (p = 1; s == 50; p = p+1+4*(p>=10)) {
    for (s = f = 0; (s < 50) and (f < 51); ) {
      invoke generator for (p,q);
      if (generator found tuple) s = s+1;
      else f = f+1;
    }
    report average times for this (p,q);
  }
}
```

We can think of the experiment as moving in a table where the rows correspond to values of p and the columns correspond to values of q . (This table underlies the reporting of the data in Appendix A.) We start at $q = n - 1$, and move down in the table (fixed q , increasing p). If, for a given p, q pair, we get 50 successes with a 50% failure rate or less, then we move down in the current column (we increase p). When a greater than 50% failure rate occurs, the current column is completed; we move left (decrease q) and start

at the top of the new column (set p to 1). The table is completed when we finish the $q = 2$ column, or stop at row 1 in some column.

The five dichotomies used for the experiments are listed below.

1. $c-e$.
2. $e-c$.
3. e_5-c_5 .
4. $acenou-sxz$.
5. $c-n$.

These dichotomies were chosen to represent a spectrum of problem types. Specifically, they were chosen for the following reasons. The $c-e$ dichotomy is believed to be a difficult one to find tuples for. The dual of this dichotomy, $e-c$, was included because it seemed useful to see how the experiment results change when the positive and negative classes are switched. The experimental results show that finding tuples for a given dichotomy may be much easier or harder than for the dichotomy's dual, at least with the generators used here. The e_5-c_5 dichotomy is different from the other four dichotomies because it consists of fifth generation photocopies. The $acenou-sxz$ dichotomy has several characters in each class. The $c-n$ dichotomy is relatively easy to find tuples for.

The experiments used 8-point Times Roman characters scanned at 300 dots per inch. The characters are shown in Figure 1. The fifth generation photocopies c_5 and e_5 are shown in the bottom row of the figure. All characters were trimmed to the smallest rectangular bounding frame before being presented to the generators.

For a given dichotomy, the generators drew tuples from a rectangular region of pixels Π . The height (width) of Π is the smallest character height (width), taken over the positive exemplars of the dichotomy. For the purpose of determining the p and q values of a given tuple relative to a given character c , the character is considered to be embedded in an infinite sea of white pixels. Effectively, the tuple is tested in all shift offsets that place some part of Π over some part of c 's smallest (inclusive) bounding box. Rotations are

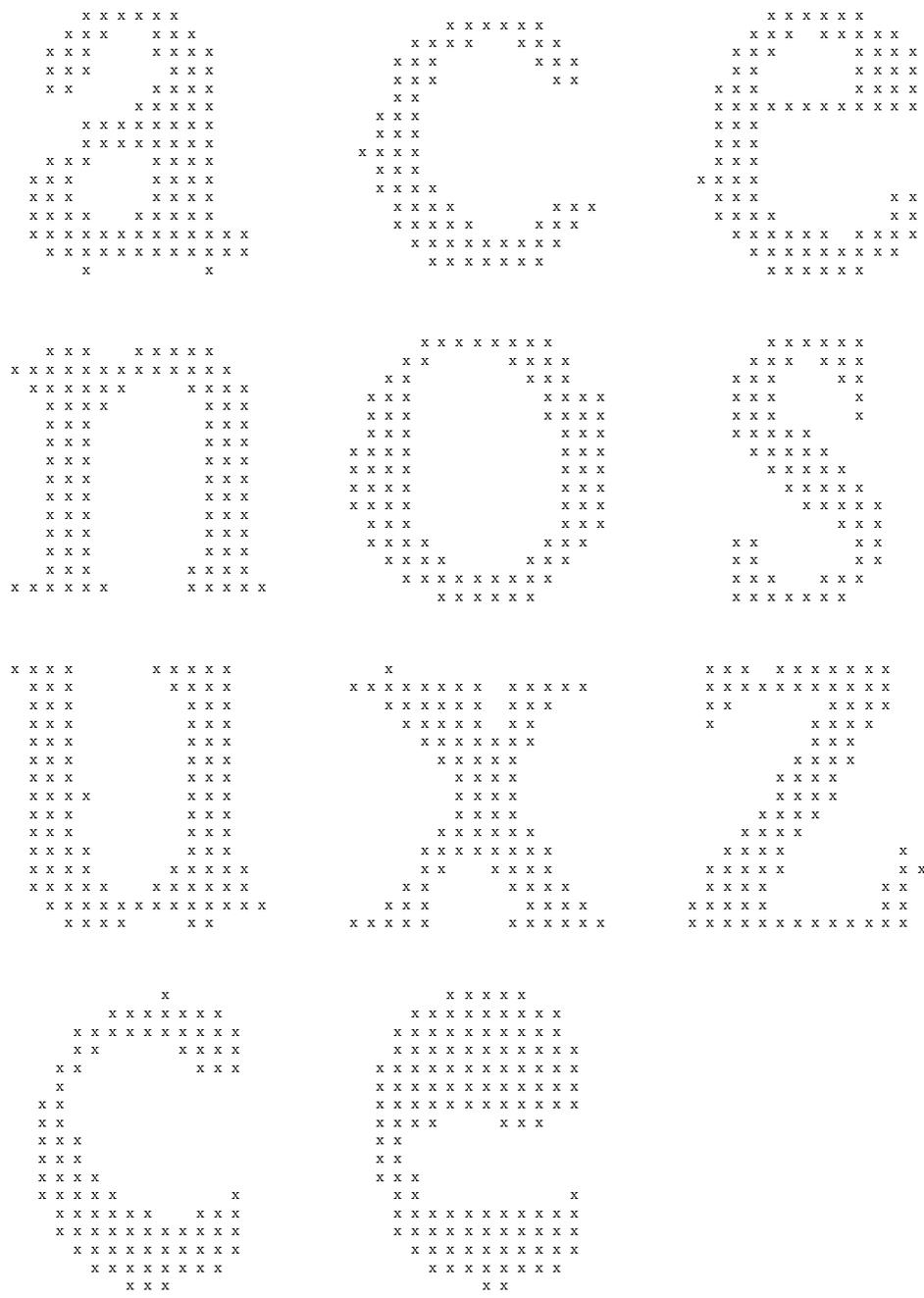


Figure 1: Characters used in experiments.

not considered, i.e., the tuple shifts are aligned with the two coordinate axes of the character.

The maximum node limit for Gen0 was 1,500,000; for Gen1, 70,000. The timeslice width for Gen1 was 500, and the width for restricted backtracking was 20. The parameters 500 and 20 were selected because of an empirically obtained belief that they give “reasonable” performance for many problems.

We did not know in advance how the generators would perform with the selected parameters on any particular problem in the experiments. The parameters were selected without regard for the particular problems included in the experiments.

In addition to the UNIX command-line arguments for the search limits given above, the following arguments were also used to invoke the generators.

```
-postkind pi -negtkind pi -randseed -one -bptime -bw
```

The `postkind pi` and `negtkind pi` arguments specify that Π and the tuple shifts are as described earlier. The `randseed` argument means that each time the generator is invoked, it uses a new seed for producing pseudorandom numbers. Thus, on successive executions, the generator may follow different execution paths. Executing the generator many times gives a sample of the generator’s possible behavior over the possible random number sequences. The `bw` argument indicates the use of black and white tuples (as opposed to all-black or all-white tuples). The `one` and `bptime` arguments are instructions related to backtracking; their precise meaning is unimportant here.

3 Description of Experiment Data

The data is presented in Appendix A. There is one table for each (dichotomy, n , generator) triple. All times are in CPU seconds.

It is convenient to explain the data by example, so we refer to Table 1. At the top of the table is the mean execution time for the $p = 1, q = n - 1$ case; here, .2132 CPU seconds. The $p = 1, q = n - 1$ case is typically the easiest (p, q) pair to find tuples for. The times within the table are scaled to this case, to clarify the changing difficulty of finding tuples as we move away from $p = 1, q = n - 1$.

Within the table, the nonblank table entries can be categorized into three types (not including shading). The first type of entry has two numbers stacked above each other, as with $p = 1, q = 2$. The $p = 1, q = 2$ entry indicates that no failures occurred for $p = 1, q = 2$; the average scaled time of the 50 successes was 103 (103 times the base time of .2132 seconds), and the standard deviation of the scaled time was 73. Another type of table entry has five numbers. For example, with $p = 10, q = 3$, the mean scaled time for the 50 successes was 544, with a standard deviation of 395. There were 42 failures; the average scaled time for these 42 failures was 1313; and the standard deviation was 37. A third type of table entry occurs when there are no successes, as with $p = 2, q = 2$. For $p = 2, q = 2$ there were 51 failures; the average failure time was 1074, with a standard deviation of 24.

When available, the tables show the optimal (largest) p values for a given value of q . Optimal values are indicated by dark shading in the appropriate table entry. (No trials were conducted for certain optimal (p, q) pairs; this happens when the generators did not achieve optimality, or when the optimal p value is larger than 10 and is not a multiple of 5.) When it is known that no solution exists for a given value of q , this is indicated by an asterisk next to the column header. The optimality of the displayed values was determined by exhaustive search. In some cases, when it was not possible to determine the optimal values, the best known solution values are indicated with light shading. For example, see Table 2.

Considerable CPU time was used in the attempt to determine optimal values. In some cases, weeks of CPU time were spent determining whether a solution exists for a given (p, q) pair.

The procedure for determining optimal values was independent of the C-like pseudocode in Section 2. This procedure followed the staircase pattern described in [1]. For example, in Table 1, to find the optimal values, the generator traversed (p, q) pairs in the following order: $(1, 2)$, $(2, 2)$, $(2, 3)$, $(3, 3)$, $(4, 3)$, $(5, 3)$, \dots , $(11, 3)$, $(12, 3)$. The optimal pairs are therefore $(1, 2)$ and $(11, 3)$, since no solutions were found for $(2, 2)$ or $(12, 3)$.

The tables were produced by a *perl* script that scans the 50 to 101 generator output files for each relevant (p, q) pair in a given table, and automatically produces the L^AT_EX source to make the table.

4 Discussion

In this section we look at the three questions that the experiments were designed to answer. In Section 4.4, we consider the quality of solution tuples as a function of the dichotomy.

4.1 Tradeoff Between Difficulty and Execution Time

Apparently, solution tuples are relatively abundant when p, q is close to $p = 1, q = n - 1$. As we move away from this point, ostensibly the problem becomes increasingly constrained and the solution tuple density decreases. For a given dichotomy and value of n , there is a tradeoff between execution time and the difficulty of the problem as defined by the (supposed) solution density. Inspection of the tables shows that Gen0 and Gen1 require more execution time as we move away from $p = 1, q = n - 1$. Thus, we can control the amount of time to find tuples by selecting p and q .

4.2 Time to Find Tuples

Here we address the issue of whether or not there exists an algorithm that finds distinguishing tuples reasonably quickly, when such tuples exist for a given dichotomy and value of n . Inspection of the tables in Appendix A yields an affirmative answer, at least for the problems in these experiments.

When solutions are relatively abundant or when solutions exist in a constrained situation such as for small n , either Gen0 or Gen1 is sufficient to find tuples. For the other cases, Gen0 does not do as well as Gen1.

4.3 Gen1 vs. Gen0

The third purpose of the experiments was to show that Gen1 is a “good” algorithm by comparing it to the benchmark algorithm Gen0.

For the most part, the measure of an algorithm should be how it does on difficult problems. (It is easy to do well on easy problems.) The definition of a difficult problem depends in general on the algorithm used to solve it. Here, we say that the difficulty increases when n increases and when (p, q) moves away from $p = 1, q = n - 1$. One reason for this is that (empirically) these cases take more time; in every table in Appendix A, as we move away from $p = 1, q = n - 1$ we see increased time to find a solution. Another reason is that the tuples away from $p = 1, q = n - 1$ are the ones that are most desirable for OCR [1].

Examination of the tables shows that for the difficult problems as defined above, Gen1 finds solutions faster and over a wider range than Gen0. For the “easy” problems, Gen0 does better.

It is instructive to examine why Gen0 does better than Gen1 on easy problems. One reason is that Gen1 was designed with difficult problems in mind. A consequence of this design is that Gen1 cannot find any solution without evaluating many tuples at each search node, even when the

search goes directly to a solution with no backtracking. On the other hand, Gen0 only evaluates one tuple at each search node. These facts and the experimental data suggest that we might want to form a hybrid algorithm that executes Gen0 and Gen1 in parallel. If t_0 and t_1 are the respective execution times for Gen0 and Gen1 on a given problem, then the hybrid algorithm takes time that is at worst roughly $2 \min(t_0, t_1)$. This is in some ways more desirable than the execution characteristics of either generator. For easy instances, the hybrid uses at most twice the time of Gen0; for difficult instances, the hybrid uses at most twice the time of Gen1.

There are easy instances in these experiments where Gen0 finds tuples over a wider range than Gen1 does, e.g., the $c-e$ dichotomy for $n = 4$. For these instances, the failure times for Gen0 are on the order of 15 or 20 minutes, whereas the failure times for Gen1 are on the order of a few seconds. This suggests that for these instances the search width parameter of 20 used for Gen1 is too restrictive; this value causes Gen1’s search tree to be small and to contain no solutions. We suspect that Gen1 would find tuples for these instances if the search width were increased. The experiments appear, in these instances, to be unfair to Gen1, since Gen0 is allowed to spend more time searching.

4.4 Quality of Tuples as a Function of the Dichotomy

Here we briefly consider how, if at all, the quality of solution tuples depends on the dichotomy.

The best known values of p for a given value of q are summarized in Figure 2 ($n = 4$), Figure 3 ($n = 7$), and Figure 4 ($n = 10$), for each of the five dichotomies. The figures represent the shaded boxes that appear in the tables in Appendix A. Figure 2 ($n = 4$) represents optimal values; Figures 3 and 4 contain some entries that may not be optimal.

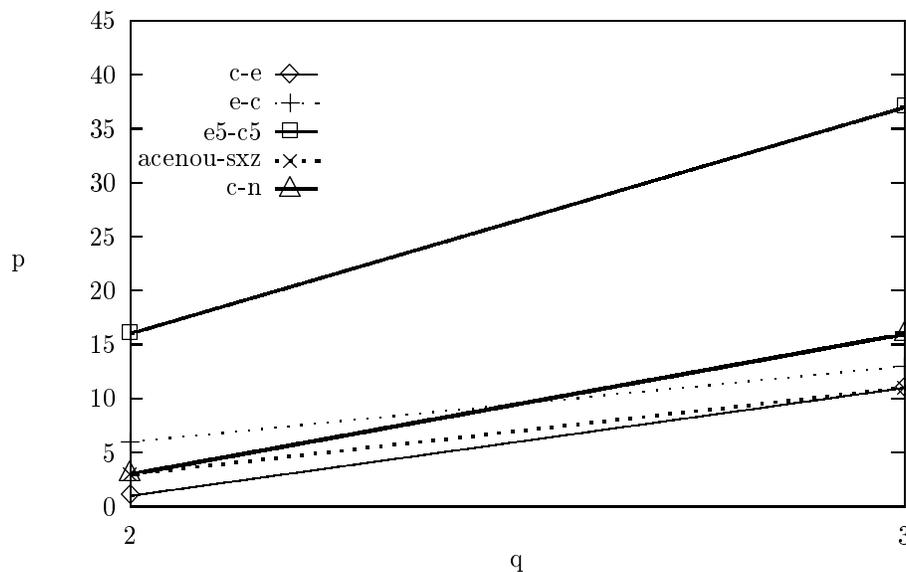


Figure 2: Optimal p for given q ($n = 4$).

A criterion for comparing the quality of tuples available to two dichotomies A and B is as follows: we can consider A to have higher quality tuples available than B if, in a plot of optimal values such as Figure 2, the curve for A is everywhere at or above the curve for B .

According to this criterion, the e_5-c_5 dichotomy has better tuples available than the other four dichotomies. Using Figures 2, 3, and 4, the five dichotomies can be roughly ordered in increasing order of quality of available tuples as follows:

$$c-e < acenou-sxz < e-c \leq c-n < e_5-c_5.$$

This ordering is consistent with our earlier belief that the $c-e$ dichotomy is a difficult one to find tuples for. The plots suggest that higher quality tuples are available for the $e-c$ dichotomy than for its converse (except when $q = n - 1$).

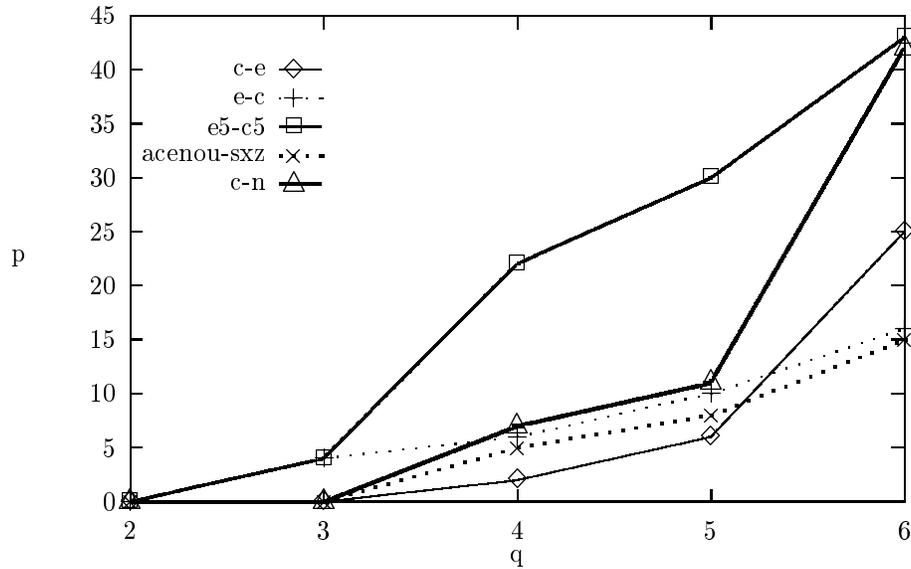


Figure 3: Largest known p for given q ($n = 7$).

The change in behavior at $q = n - 1$ is unexplained. It may be an artifact of the fact that the plots use only the best known values of p , and not necessarily the optimal values. This is another reminder that the data reported in this section is not exact and should be treated cautiously.

Despite their non-exact nature, the plots make it fairly clear that the quality of available tuples varies widely across dichotomies.

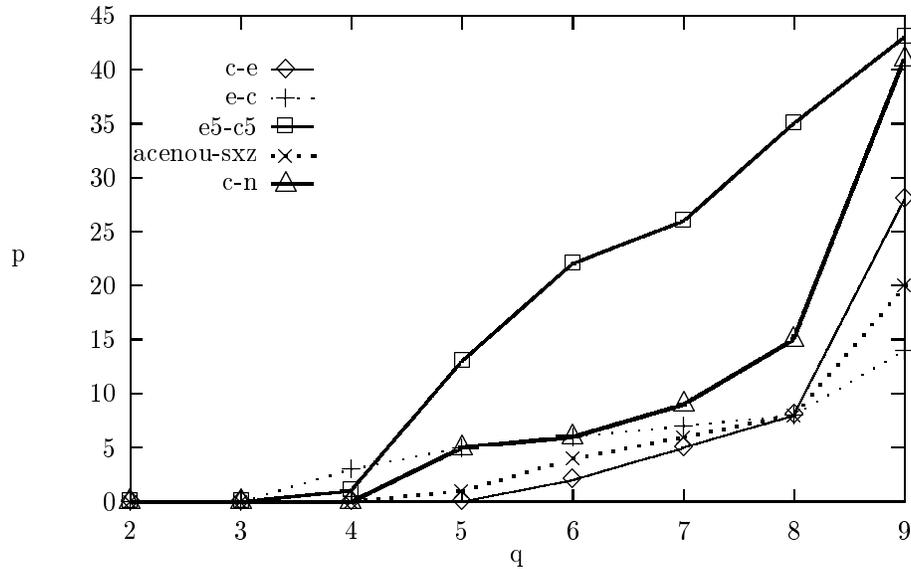


Figure 4: Largest known p for given q ($n = 10$).

5 Acknowledgement

I would like to thank D. Jung, M. Krishnamoorthy, and G. Nagy for commenting on a preliminary version of this report. G. Nagy participated in the design of the experiments, and suggested the plots in Section 4.4. Scanned characters were supplied by D. Jung.

References

- [1] D. Jung and M. Krishnamoorthy and G. Nagy and A. Shapira, “N-Tuple Features for OCR Revisited,” Submitted to IEEE Transactions on Pattern Analysis and Machine Intelligence, 1995.

Appendix A: Experiment Data

This appendix contains 30 tables, one for each (dichotomy, n , generator) triple.

Note: to reduce the table width, Table 21 does not have a column for $q = 3$. The data that would go in this column is that for $p = 1, q = 3$ there were 51 failures with an average scaled time of 140 and a standard deviation of 2.5. Exhaustive search determined that there are no solutions for $q = 3$.

Mean time for $p=1, q=3$: .2132 s

	q	
	2	3
1	103 73	1 0.08
2	1074 ⁵¹ 24	1.0 0.17
3		1.6 1.2
4		2.6 3.2
5		4.1 3.8
6		7.6 7.8
7		23 22
8		149 230
9		182 255
p		
10		544 1313 ⁴² 395 37
11		
15		1296 ⁵¹ 40
20		
25		
30		
35		
40		

Table 1: Success and failure times, $n = 4$, c-e, Gen0.

Mean time for $p=1, q=6$: .169 s

	q						
	2*	3*	4		5	6	
1		1390 ⁵¹ 89	327 345	1328 ³³ 80	2.7 3.4	1 0.19	
2			1338 0	1330 ⁵¹ 132	57 119	1.1 0.26	
3				185 246	1495 ⁵ 206	1.3 0.77	
4				375 382	1366 ²⁶ 154	8.5 30	1634 ¹ 0
5				467 482	1403 ⁵¹ 174	49 135	1469 ⁴ 142
6						43 75	1083 ⁴ 87
7						100 223	1153 ¹¹ 143
8						172 276	1230 ⁹ 247
9						186 281	1211 ³⁰ 251
p						298 281	1314 ⁴⁸ 229
10						273 152	1246 ⁵¹ 198
15							
20							
22							
25							
30							
35							
40							

Table 2: Success and failure times, $n = 7$, c-e, Gen0.

Mean time for $p=1, q=9$: .2066 s

	q									
	4	5	6		7		8		9	
1			442	1283 ⁵¹	113	1354 ³	2.2		1	
			352	186	261	106	4.4		0.097	
2					307	1347 ⁴⁸	16	1049 ¹	1.0	
					336	293	51	0	0.09	
3					448	1345 ⁵¹	130	1509 ²⁶	3.7	1652 ⁴
					364	281	214	334	15	107
4							260	1296 ⁵¹	43	1283 ⁵
							343	306	263	325
5									68	1838 ⁹
									229	149
6									125	1396 ¹⁶
									321	331
7									125	1388 ²⁸
									312	324
8									163	1202 ³²
									318	302
9									107	1375 ⁵¹
p									190	350
10										
15										
20										
25										
28										
30										
35										
40										

Table 3: Success and failure times, $n = 10$, c-e, Gen0.

Mean time for $p=1, q=3$: .1712 s

	q	
	2	3
1	6.9 6.9	1 0.13
2	8.1 6.3	0.99 0.13
3	8.3 5.8	1.0 0.16
4	6.7 5.5	0.98 0.17
5	7.6 5.2	1.0 0.16
6	37 39	1.1 0.16
7	1190 ⁵¹ 38	1.1 0.22
8		1.5 1.0
9		1.9 1.4
p		4.7 4.2
10		
13		
15		1362 ⁵¹ 52
20		
25		
30		
35		
40		

Table 4: Success and failure times, $n = 4$, e-c, Gen0.

Mean time for $p=1, q=6$: .2062 s

	q								
	2*	3		4	5	6			
1	1277 ⁵¹ 19	174 238	1278 ¹⁰ 22	2.6 2.2	1.0 0.13	1 0.033			
2		186 201	1228 ⁴² 30	18 28	1.9 2.8	4.3 23			
3		413 306	1176 ⁵¹ 73	60 108	1215 ¹ 0	3.4 11	1.1 0.24		
4				192 294	1302 ⁴ 18	12 25	1690 ¹ 0	6.3 23	
5				333 385	1242 ²³ 123	23 48	1458 ¹ 0	6.4 34	1124 ¹ 0
6				418 411	1249 ⁵¹ 132	175 266	1057 ⁶ 119	27 123	
7						378 362	1273 ⁴¹ 184	38 106	1357 ³ 142
8						723 396	1313 ⁵¹ 190	42 93	1070 ⁶ 218
9								145 279	1173 ¹⁴ 243
p								212 329	1180 ²⁰ 285
15									1205 ⁵¹ 277
16									
20									
25									
30									
35									
40									

Table 5: Success and failure times, $n = 7$, e-c, Gen0.

Mean time for $p=1, q=9$: .2076 s

	q											
	4		5		6		7		8		9	
1	290	1303 ⁵¹	184	1242 ⁶	11		13		1.0		1	
	364	74	225	122	24		83		0.16		0.04	
2			397	1182 ⁵¹	109	1250 ⁵	23	1078 ²	20		1.0	
			349	95	225	205	123	180	98		0.066	
3					248	1232 ²⁰	79	1145 ⁷	22	1548 ⁹	3.9	
					223	177	175	203	82	246	15	
4					415	1244 ⁵¹	217	1378 ²¹	35	1328 ¹⁴	11	1070 ¹
					374	215	309	260	94	268	33	0
5							306	1254 ⁵¹	102	1207 ⁴⁰	14	1435 ⁷
							274	259	207	300	44	345
6									191	1276 ⁵¹	82	1231 ²²
									230	276	209	380
7											144	1232 ³⁶
											263	379
8											198	1156 ⁵¹
											267	310
9												
p												
10												
14												
15												
20												
25												
30												
35												
40												

Table 6: Success and failure times, $n = 10$, e-c, Gen0.

Mean time for $p=1, q=3$: .2198 s

	q	
	2	3
1	1.1 0.13	1 0.044
2	1.2 0.23	0.98 0.083
3	1.3 0.29	1.0 0.054
4	1.5 0.48	1.0 0.06
5	1.8 0.65	1.0 0.057
6	2.1 1.3	1.0 0.073
7	2.1 1.1	0.99 0.083
8	2.2 1.3	1.0 0.12
9	3.1 1.8	0.99 0.071
p		
10	3.2 1.9	1.0 0.08
15	9.8 15	1.3 0.8
16		
20	1140 ⁵¹ 8.9	1.2 0.6
25		1.8 1.3
30		4.2 3.1
35		371 1263 ¹⁸ 391 30
37		
40		1223 ⁵¹ 48

Table 7: Success and failure times, $n = 4$, e_5-c_5 , Gen0.

Mean time for $p=1, q=6$: .1924 s

	q					
	2^*	3		4	5	
1	1299 ⁵¹ 47	99 106	2.0 1.6	0.99 0.19	1 0.19	
2		307 1156 ⁷ 300 76	4.9 11	1.0 0.2	0.99 0.18	
3		511 1119 ⁵¹ 335 58	16 32	1.1 0.34	1.0 0.32	
4			9.6 16	1.5 2.7	1.0 0.24	
5			14 25	25 164	1.0 0.36	
6			15 25	1.2 0.58	2.6 11	
7			22 38	2.3 4.1	3.0 1268 ¹ 9.3 0	
8			34 64	35 124	5.3 24	
9			59 83	22 74	1.6 1320 ¹ 2.9 0	
10			59 1071 ² 109 93	3.7 910 ¹ 5.4 0	5.3 1118 ² 18 179	
p			278 1075 ¹⁰ 285 150	55 913 ¹ 171 0	65 1008 ⁴ 154 93	
15			496 1047 ⁵⁰ 328 141	100 1021 ⁹ 166 220	53 1038 ¹⁰ 136 178	
20						
22						
25			1070 ⁵¹ 154	354 1086 ⁵¹ 254 189	83 1072 ¹³ 202 194	
30					161 1067 ⁵¹ 213 220	
35						
40						
43						
45						

Table 8: Success and failure times, $n = 7$, e_5-c_5 , Gen0.

Mean time for $p=1, q=9$: .2212 s

	q											
	4		5		6		7		8		9	
1	500	1333 ⁵¹	214	1313 ¹²	28		24		1.0		1	
	254	75	326	123	131		139		0.063		0.065	
2			314	1291 ²³	46	1447 ¹	22		2.9		1.0	
			362	118	146	0	116		10		0.086	
3			410	1197 ³⁸	63	1197 ²	8.4	1527 ²	23		1.1	
			342	133	159	157	36	191	120		0.6	
4			559	1170 ⁵¹	155	1328 ³	15	1719 ²	13	1198 ⁴	1.1	
			448	101	280	95	65	63	55	389	0.47	
5					68	1151 ⁹	26	1491 ¹	16	1485 ⁴	1.3	1027 ⁴
					147	199	98	0	91	245	1.2	190
6					132	1136 ⁸	28	1225 ⁵	40	1194 ³	1.7	1267 ³
					208	236	68	318	151	201	4.0	523
7					198	1119 ¹⁹	33	1256 ⁷	48	1287 ¹⁰	5.6	1343 ³
					254	156	65	314	150	253	32	285
8					222	1102 ¹²	36	1123 ¹¹	45	1080 ⁷	2.3	999 ⁶
					240	174	76	203	147	242	5.2	235
9					168	1050 ⁴²	105	1115 ¹²	53	1168 ⁸	8.4	1264 ⁶
					190	159	271	261	116	226	30	449
10					279	1091 ³¹	67	1082 ⁷	50	1185 ¹⁷	42	1376 ⁹
					269	174	128	256	138	290	140	348
13												
p												
15					181	990 ⁵¹	112	971 ⁵¹	70	1244 ²¹	18	1061 ¹⁸
					195	150	158	164	166	343	74	301
20									131	1154 ⁵¹	78	1095 ⁴⁷
									245	302	193	277
22												
25											59	1014 ⁵¹
											121	232
26												
30												
35												
40												
43												
45												

Table 9: Success and failure times, $n = 10$, e_5 - c_5 , Gen0.

Mean time for $p=1, q=3$: .7872 s

	q	
	2	3
1	6.2 5.7	1 0.28
2	11 12	1.3 0.67
3	23 23	1.4 0.72
4	564^{51} 7.4	1.8 1.1
5		3.4 3.7
6		7.5 15
7		12 31
8		21 44
9		58 110
p		137 185
10		835^5 26
11		
15		822^{51} 22
20		
25		
30		
35		
40		

Table 10: Success and failure times, $n = 4$, acenou-sxz, Gen0.

Mean time for $p=1, q=6$: 1.9652 s

	q							
	2^*	3^*	4		5		6	
1		324^{51} 42	136 121	345^{14} 27	9.7 23	272^2 14	1 3.0	
2			131 121	348^{21} 27	33 51	249^8 31	2.3 6.3	
3			138 114	338^{51} 34	64 80	279^{22} 70	6.3 236^4 15 57	
4					89 98	282^{38} 68	18 46	236^9 53
5					114 115	313^{51} 69	25 48	276^{13} 53
6							35 53	252^{27} 47
7							51 72	263^{51} 50
8								
p								
9								
10								
15								
20								
25								
30								
35								
40								

Table 11: Success and failure times, $n = 7$, acenou-sxz, Gen0.

Mean time for $p=1, q=9$: 2.1424 s

	q								
	4	5	6	7		8		9	
1				43 58	216 ⁵¹ 40	13 34	193 ¹⁴ 36	1 2.6	169 ¹ 0
2						39 52	217 ⁴¹ 42	6.6 18	219 ³ 68
3						57 67	203 ⁵¹ 45	14 33	186 ¹⁸ 50
4								18 37	180 ³⁸ 46
5								30 40	177 ⁵¹ 43
6									
7									
8									
p									
9									
10									
15									
20									
25									
30									
35									
40									

Table 12: Success and failure times, $n = 10$, acenou-sxz, Gen0.

Mean time for $p=1, q=3$: .1804 s

	q	
	2	3
1	12 8.6	1 0.17
2	31 35	0.98 0.17
3	63 62	1.0 0.2
4	1297 ⁵¹ 36	0.94 0.16
5		1.0 0.2
6		1.1 0.5
7		1.1 0.58
8		1.4 1.2
9		1.5 1.4
p		2.4 2.2
10		154 221
15		
16		
20		1292 ⁵¹ 90
25		
30		
35		
40		

Table 13: Success and failure times, $n = 4$, c-n, Gen0.

Mean time for $p=1, q=6$: .2064 s

	q				
	2*	3*	4	5	6
1		1199 ⁵¹ 65	21 28	1.2 0.35	1 0.077
2			93 131	1.6 1.5	1.0 0.12
3			140 1141 ²² 219 110	3.4 5.8	1.0 0.091
4			183 1200 ²⁰ 229 111	45 122	1.0 0.082
5			233 1168 ³¹ 281 109	58 1604 ¹ 132 0	2.1 6.1
6			402 1197 ⁵¹ 370 127	94 1101 ⁶ 169 148	1.3 1.0
7				114 1073 ⁸ 194 192	7.5 34
8				227 1127 ¹⁷ 303 234	8.2 30
9				430 1241 ⁵¹ 338 196	22 1114 ² 54 121
10					79 1368 ⁵ 270 284
p					
11					
15					183 1372 ¹⁴ 234 315
20					332 1208 ⁵¹ 325 287
25					
30					
35					
40					
42					
45					

Table 14: Success and failure times, $n = 7$, $c-n$, Gen0.

Mean time for $p=1, q=9$: .1706 s

	q									
	4	5		6		7		8		9
1		510	1242 ⁵¹	125	1066 ⁶	7.2	953 ¹	1.0		1
		422	136	239	131	17	0	0.28		0.16
2				170	1154 ²⁶	16	931 ¹	1.3		0.95
				242	158	61	0	1.2		0.13
3				268	1147 ⁴⁶	106	1123 ¹¹	24	859 ¹	1.0
				295	143	232	281	102	0	0.42
4				309	1186 ⁵¹	161	1231 ²⁴	52	1037 ³	1.1
				315	178	255	253	113	81	0.4
5						155	1099 ⁵¹	36	1049 ¹¹	1.7
						251	179	85	154	3.7
6								149	1177 ¹³	4.1
								249	282	17
7								184	1134 ³²	26
								302	254	117
8								128	1215 ⁵¹	34
								187	341	118
9										24
										72
p										1039 ⁷
10										36
										1199 ²³
										82
15										142
										1272 ²⁹
										247
20										96
										1297 ⁵¹
										168
25										363
30										
35										
40										
41										
45										

Table 15: Success and failure times, $n = 10$, $c-n$, Gen0.

Mean time for $p=1, q=3$: .6368 s

	q	
	2	3
1	2.4 ⁵¹ 0.044	1 0.049
2		0.95 0.055
3		0.98 0.11
4		1.3 0.43
5		4.8 ⁵¹ 0.18
6		
7		
8		
9		
p		
10		
11		
15		
20		
25		
30		
35		
40		

Table 16: Success and failure times, $n = 4$, c-e, Gen1.

Mean time for $p=1, q=6$: .5668 s

	q				
	2*	3*	4	5	6
1		15 ⁵¹ 1.8	86 64	1.1 0.2	1 0.096
2			166 ⁵¹ 7.9	7.8 3.8	0.99 0.12
3				224 352 ⁵¹ 6.3 16	2.5 1.8
4					2.3 1.4
5					1.5 0.73
6					2.3 0.62
7					3.1 1.3
8					3.5 1.9
9					3.8 2.5
p					5.4 4.1
10					94 467 ⁷ 66 27
15					457 ⁵¹ 17
20					
22					
25					
30					
35					
40					

Table 17: Success and failure times, $n = 7$, c-e, Gen1.

Mean time for p=1,q=9: .8808 s

	<i>q</i>					
	4	5	6	7	8	9
1			83 191 ⁵¹ 3.0 7.1	0.98 0.04	1.0 0.041	1 0.054
2				33 221 ⁹ 23 16	5.4 2.6	0.92 0.078
3				171 ⁵¹ 13	5.3 0.59	2.2 1.4
4					6.9 3.0	2.2 1.4
5					8.1 158 ⁸ 9.0 9.6	1.1 0.47
6					150 ⁵¹ 8.9	1.9 0.42
7						2.4 0.97
8						3.0 1.3
9						2.2 1.1
<i>p</i>						2.9 1.4
10						3.9 3.3
15						5.5 125 ² 4.0 3.5
20						109 ⁵¹ 10
25						
28						
30						
35						
40						

Table 18: Success and failure times, n = 10, c-e, Gen1.

Mean time for $p=1, q=3$: .472 s

	q	
	2	3
1	1.2 0.14	1 0.12
2	1.2 0.15	1.1 0.12
3	1.2 0.14	1.0 0.14
4	1.2 0.12	1.0 0.11
5	3.0^{51} 0.29	1.0 0.1
6		1.1 0.13
7		1.2 0.12
8		1.2 0.13
9		1.2 0.18
p		1.2 0.17
10		
13		
15		4.6^{51} 0.44
20		
25		
30		
35		
40		

Table 19: Success and failure times, $n = 4$, e-c, Gen1.

Mean time for $p=1, q=6$: .6548 s

	q				
	2*	3	4	5	6
1	4.8 ⁵¹ 0.15	1.4 0.11	1.6 0.21	1.7 0.25	1 0.2
2		1.5 0.18	1.3 0.18	1.5 0.25	1.2 0.19
3		44 7.2	1.5 0.29	1.6 0.27	1.0 0.18
4		44 ⁵¹ 1.3	1.3 0.22	1.4 0.27	1.0 0.16
5			3.0 3.3	1.4 0.26	0.97 0.13
6			266 44	2.6 1.7	1.2 0.15
7			400 ⁵¹ 2.8	6.0 3.3	1.2 0.021
8				9.5 7.3	1.2 0.03
9				57 336 ⁵¹ 5.7 7.7	1.1 0.079
p					1.0 0.029
10					
15					203 ⁵¹ 7.2
16					
20					
25					
30					
35					
40					

Table 20: Success and failure times, $n = 7$, e-c, Gen1.

Mean time for $p=1, q=9$: .7346 s

	q					
	4	5	6	7	8	9
1	4.3 2.9	1.7 0.23	1.8 0.21	1.9 0.18	1.9 0.21	1 0.29
2	87 151 ⁴⁷ 43 7.1	1.4 0.22	1.4 0.18	1.6 0.24	1.5 0.21	1.2 0.27
3	267 ⁵¹ 9.8	1.9 0.72	1.7 0.23	1.7 0.24	1.6 0.26	0.96 0.23
4		42 334 ⁴ 67 9.2	1.4 0.21	1.4 0.2	1.4 0.18	1.1 0.21
5		378 ⁵¹ 19	2.9 3.3	3.1 3.6	1.3 0.18	0.98 0.17
6			415 ⁵¹ 20	50 416 ⁵¹ 18 19	2.0 1.7	1.4 0.036
7					4.9 3.6	1.3 0.017
8					8.1 3.3	1.2 0.079
9					318 ⁵¹ 12	1.1 0.13
p						
10						1.0 0.052
14						
15						198 ⁵¹ 8.7
20						
25						
30						
35						
40						

Table 21: Success and failure times, $n = 10$, e-c, Gen1 (see note at beginning of Appendix A).

Mean time for $p=1, q=3$: .6522 s

	q	
	2	3
1	1.2 0.087	1 0.1
2	1.1 0.093	1.0 0.12
3	1.1 0.081	1.0 0.095
4	1.0 0.053	1.1 0.1
5	1.1 0.045	0.91 0.037
6	1.0 0.056	1.2 0.02
7	1.1 0.11	1.2 0.024
8	1.1 0.073	1.2 0.028
9	1.2 0.14	1.2 0.021
p		
10	1.3 0.18	1.2 0.021
15	1.4 0.1	1.1 0.021
16		
20	3.1^{51} 0.038	1.0 0.023
25		1.2 0.18
30		1.3 0.2
35		3.3^{51} 0.097
37		
40		

Table 22: Success and failure times, $n = 4$, e_5 - c_5 , Gen1.

Mean time for $p=1, q=6$: .6554 s

	q				
	2*	3	4	5	6
1	2.8 ⁵¹ 0.51	8.9 6.1	2.8 0.95	1.0 0.25	1 0.25
2		8.5 3.6	2.0 1.1	2.3 1.1	0.98 0.27
3		54 20	1.5 0.7	2.0 0.89	0.99 0.25
4		156 ⁵¹ 11	2.2 0.88	1.3 0.26	1.2 0.34
5			3.6 2.5	1.6 0.33	0.78 0.14
6			1.4 0.76	1.3 0.23	1.3 0.23
7			3.3 1.2	1.2 0.21	1.2 0.21
8			2.5 2.0	1.2 0.21	1.2 0.22
9			2.4 1.9	1.1 0.21	1.2 0.21
10			3.6 2.5	1.1 0.15	1.2 0.21
p					
15			5.3 4.4	1.9 1.2	1.1 0.19
20			5.4 1.2	2.3 1.5	0.99 0.17
22					
25			56 ⁵¹ 8.5	6.3 8.9	1.7 0.97
30				145 ⁵¹ 8.4	3.4 123 ¹ 5.0 0
35					150 ⁵¹ 8.4
40					
43					
45					

Table 23: Success and failure times, $n = 7$, e_5 - c_5 , Gen1.

Mean time for $p=1, q=9$: .8522 s

	q						
	4	5	6	7	8	9	
1	30 18 2.4	132 ⁵¹ 0.99	2.6 0.93	2.6 0.93	2.7 0.85	1.1 0.27	1 0.24
2		2.0 1.2	2.4 1.2	2.3 1.2	2.2 1.2	1.0 0.27	
3		3.7 2.8	1.7 0.75	2.0 0.94	2.0 0.99	0.94 0.19	
4		6.8 3.0	2.2 0.87	2.3 0.71	1.5 0.092	1.4 0.3	
5		11 5.3	3.5 2.2	3.2 1.7	1.9 0.2	0.77 0.022	
6		10 4.1	1.7 0.83	1.7 0.85	1.5 0.021	1.5 0.049	
7		13 7.5	3.8 1.3	3.7 1.2	1.4 0.062	1.4 0.064	
8		28 31	3.5 2.4	2.7 1.7	1.4 0.064	1.4 0.05	
9		63 239 ⁹ 49 13	3.2 2.5	2.8 1.9	1.3 0.03	1.4 0.041	
10		42 262 ⁴⁴ 19 20	4.6 3.0	4.1 2.7	1.3 0.047	1.4 0.055	
13							
p							
15		265 ⁵¹ 11	13 5.6	3.5 2.6	1.5 0.84	1.1 0.055	
20			74 216 ⁵¹ 65 12	11 8.2	2.8 1.4	0.96 0.033	
22							
25				130 ⁵¹ 15	2.6 132 ³³ 1.1 16	1.7 0.81	
26							
30					118 ⁵¹ 5.5	3.0 126 ⁷ 3.5 6.0	
35						144 ⁵¹ 6.8	
40							
43							
45							

Table 24: Success and failure times, $n = 10$, e_5 - c_5 , Gen1.

Mean time for $p=1, q=3$: 1.6832 s

	q	
	2	3
1	1.0 0.18	1 0.15
2	1.1 0.2	0.93 0.093
3	1.1 0.24	0.89 0.099
4	2.5^{51} 0.4	0.89 0.12
5		0.95 0.14
6		0.91 0.15
7		0.96 0.15
8		0.9 0.15
9		0.9 0.16
p		
10		0.96 0.18
11		
15		5.4^{51} 0.75
20		
25		
30		
35		
40		

Table 25: Success and failure times, $n = 4$, acenou-sxz, Gen1.

Mean time for $p=1, q=6$: 2.235 s

	q				
	2^*	3^*	4	5	6
1		64^{51} 3.5	1.2 0.32	1.0 0.012	1 0.032
2			2.5 1.5	1.4 0.8	0.95 0.019
3			6.5 5.2	0.99 0.26	0.91 0.011
4			13 5.4	1.4 0.59	0.88 0.0089
5			12 3.0	2.0 1.3	1.8 1.2
6			110^{51} 1.5	5.1 6.5	0.87 0.028
7				135^{51} 2.9	0.87 0.0054
8					0.83 0.01
p					0.81 0.008
9					0.97 0.3
10					
15					422^{51} 6.1
20					
25					
30					
35					
40					

Table 26: Success and failure times, $n = 7$, acenou-sxz, Gen1.

Mean time for $p=1, q=9$: 2.216 s

	q					
	4	5	6	7	8	9
1		196 ⁵¹ 9.5	70 32	3.0 2.1	1.0 0.16	1 0.17
2			47 182 ⁵ 39 2.9	3.0 3.6	1.3 0.77	0.88 0.13
3			87 173 ¹² 45 11	14 9.0	1.1 1.0	0.89 0.14
4			186 ⁵¹ 19	9.3 7.3	1.7 0.56	0.84 0.14
5				13 3.6	1.8 1.1	1.5 1.2
6				139 ⁵¹ 6.1	6.0 135 ² 7.9 4.0	0.82 0.14
7					113 ⁵¹ 12	0.84 0.13
8						0.79 0.13
9						0.75 0.13
10						0.88 0.3
15						264 ⁵¹ 17
20						
25						
30						
35						
40						

Table 27: Success and failure times, $n = 10$, acenou-sxz, Gen1.

Mean time for $p=1, q=3$: .5378 s

	q	
	2	3
1	1.9 0.38	1 0.078
2	2.8 ⁵¹ 0.44	0.98 0.084
3		0.98 0.092
4		0.98 0.086
5		0.95 0.093
6		0.95 0.089
7		0.95 0.13
8		1.1 0.16
9		1.0 0.18
p		1.0 0.18
10		1.0 0.18
15		0.99 0.21
16		
20		6.4 ⁵¹ 1.1
25		
30		
35		
40		

Table 28: Success and failure times, $n = 4$, c-n, Gen1.

Mean time for $p=1, q=6$: .8056 s

	q				
	2*	3*	4	5	6
1		75 ⁵¹ 3.0	1.8 1.4	1.0 0.024	1 0.042
2			3.9 3.2	1.0 0.019	0.99 0.04
3			14 4.1	0.95 0.016	0.95 0.035
4			23 12	4.4 4.3	0.95 0.044
5			47 55	13 6.4	0.97 0.024
6			351 ⁵¹ 2.7	15 6.7	0.93 0.034
7				11 5.4	0.93 0.034
8				16 4.7	0.91 0.02
9				16 330 ¹⁴ 6.3 4.4	0.92 0.018
10				312 ⁵¹ 6.5	0.89 0.021
p					
11					
15					2.7 1.8
20					13 6.3
25					100 92
30					429 535 ¹³ 99 11
35					232 414 ⁵¹ 161 13
40					
42					
45					

Table 29: Success and failure times, $n = 7$, $c-n$, Gen1.

Mean time for $p=1, q=9$: .8262 s

	q					
	4	5	6	7	8	9
1	173 ⁵¹ 26	70 144 ²⁸ 49 11	1.9 1.4	1.1 0.63	0.97 0.17	1 0.065
2		182 ⁵¹ 24	10 3.7	1.1 0.54	1.0 0.18	0.94 0.046
3			16 3.7	5.5 4.3	0.99 0.17	1.0 0.14
4			62 266 ¹³ 67 14	8.1 5.0	1.2 0.48	1.1 0.12
5			104 328 ⁴ 110 13	16 6.6	13 8.4	1.0 0.19
6			335 ⁵¹ 13	19 4.3	17 7.0	0.9 0.16
7				357 ⁵¹ 10	14 6.2	0.92 0.17
8					17 6.4	0.88 0.16
9					18 365 ²² 7.2 13	0.94 0.17
p					342 ⁵¹ 16	0.91 0.16
10						2.2 1.6
15						23 14
20						91 494 ²⁶ 60 20
25						346 435 ⁵¹ 91 27
30						
35						
40						
41						
45						

Table 30: Success and failure times, $n = 10$, $c-n$, Gen1.