

Experiments on the Generation of Distinguishing N-Tuples for Selected Character Dichotomies

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1 Introduction

This report is a companion to [1]. In the present document, we describe experiments on generating n-tuples for optical character recognition. The focus is on the generation process itself, and not the use of the generated tuples. The experiments were conducted at Rensselaer during late January, 1995.

An outline of the report is as follows. In Section 2 we describe the experiments. The presentation of the experiment data is explained in Section 3. Section 4 contains some discussion. The experiment data is presented in Appendix A.

This report assumes familiarity with [1]. Additional information regarding the generators, and further discussion of the experiments, may be found there.

2 Description of Experiments

We consider the generation of n -tuples using two generators, called Gen0 and Gen1. Gen0 is a simple backtracking algorithm. Gen1 is a more sophisticated backtracking algorithm.

The main purpose of the experiments was to show the following.

1. Distinguishing tuples can be found reasonably quickly, when they exist.
2. Gen1 is a reasonably efficient algorithm for generating tuples. This was to be shown by comparing Gen1 to the benchmark program Gen0.
3. The execution time of Gen1 can be controlled by varying the difficulty of the problem as defined by the p and q parameters.

We also consider the quality of solution tuples as a function of the dichotomy.

The experiments consisted of executing Gen0 and Gen1 on selected character dichotomies for $n = 4, 7, 10$. Five dichotomies were used, giving $2 \cdot 3 \cdot 5 = 30$ experiments. For each experiment, Appendix A contains one table summarizing the results. Combined, the experiments represent more than 40,000 generator invocations, and hundreds of hours of CPU time. (Unless otherwise specified, all times in this document are CPU times.)

The experiments were conducted on 4 two-processor SPARC 20's, using one CPU at a time. Three of the SPARC 20's had 64 megabytes of memory, and one had 256 megabytes. The generation process had modest memory requirements, e.g., a megabyte or less, so the amount of memory possessed by the machines was not critical for these experiments.

The generators were executed in the mode "find a solution for a specified (p, q) pair." The generators executed until either a solution was found, or the search terminated because a maximum search node limit was reached. The maximum node limits were selected so that the generators executed for

about 5 minutes on a certain benchmark problem, when no solution was found. This “failure time” is different for different problems, because the time spent at a given search node varies. For these experiments, failure times were usually between 2 and 25 minutes. In some cases, when the problem was extremely constrained as when q is very small, failure times were only a few seconds.

For a given dichotomy, value of n , and generator, the generator was executed according to the following C-like pseudocode.

```
integer p,q,s,f;    // for given p,q: s=successes, f=failures

for (q = n-1, p = 0, s = 50; (q >= 2) and (p != 1); q = q-1) {
  for (p = 1; s == 50; p = p+1+4*(p>=10)) {
    for (s = f = 0; (s < 50) and (f < 51); ) {
      invoke generator for (p,q);
      if (generator found tuple) s = s+1;
      else f = f+1;
    }
    report average times for this (p,q);
  }
}
```

We can think of the experiment as moving in a table where the rows correspond to values of p and the columns correspond to values of q . (This table underlies the reporting of the data in Appendix A.) We start at $q = n - 1$, and move down in the table (fixed q , increasing p). If, for a given p, q pair, we get 50 successes with a 50% failure rate or less, then we move down in the current column (we increase p). When a greater than 50% failure rate occurs, the current column is completed; we move left (decrease q) and start

at the top of the new column (set p to 1). The table is completed when we finish the $q = 2$ column, or stop at row 1 in some column.

The five dichotomies used for the experiments are listed below.

1. $c-e$.
2. $e-c$.
3. e_5-c_5 .
4. $acenou-sxz$.
5. $c-n$.

These dichotomies were chosen to represent a spectrum of problem types. Specifically, they were chosen for the following reasons. The $c-e$ dichotomy is believed to be a difficult one to find tuples for. The dual of this dichotomy, $e-c$, was included because it seemed useful to see how the experiment results change when the positive and negative classes are switched. The experimental results show that finding tuples for a given dichotomy may be much easier or harder than for the dichotomy's dual, at least with the generators used here. The e_5-c_5 dichotomy is different from the other four dichotomies because it consists of fifth generation photocopies. The $acenou-sxz$ dichotomy has several characters in each class. The $c-n$ dichotomy is relatively easy to find tuples for.

The experiments used 8-point Times Roman characters scanned at 300 dots per inch. The characters are shown in Figure 1. The fifth generation photocopies c_5 and e_5 are shown in the bottom row of the figure. All characters were trimmed to the smallest rectangular bounding frame before being presented to the generators.

For a given dichotomy, the generators drew tuples from a rectangular region of pixels Π . The height (width) of Π is the smallest character height (width), taken over the positive exemplars of the dichotomy. For the purpose of determining the p and q values of a given tuple relative to a given character c , the character is considered to be embedded in an infinite sea of white pixels. Effectively, the tuple is tested in all shift offsets that place some part of Π over some part of c 's smallest (inclusive) bounding box. Rotations are

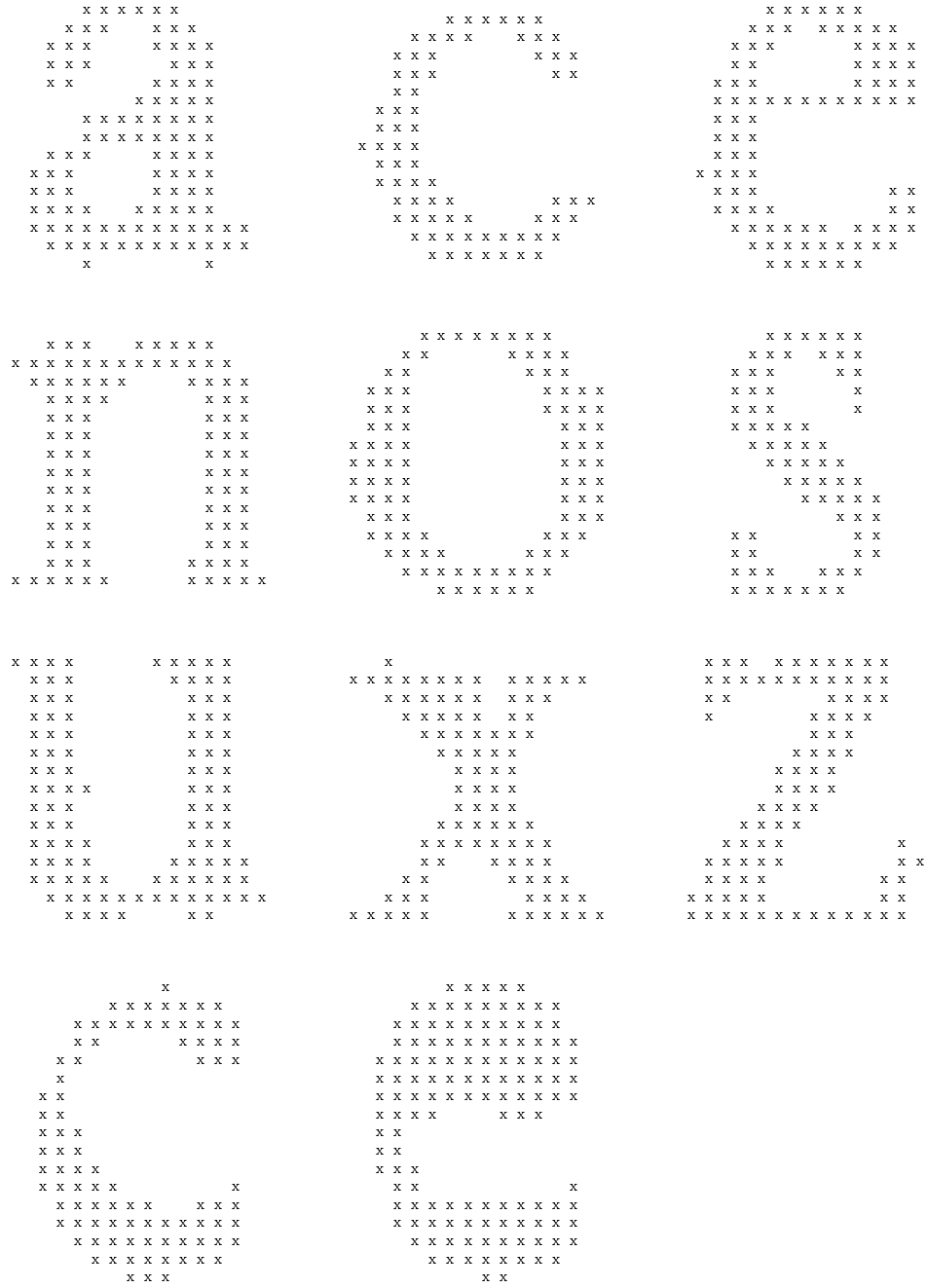


Figure 1: Characters used in experiments.

not considered, i.e., the tuple shifts are aligned with the two coordinate axes of the character.

The maximum node limit for Gen0 was 1,500,000; for Gen1, 70,000. The timeslice width for Gen1 was 500, and the width for restricted backtracking was 20. The parameters 500 and 20 were selected because of an empirically obtained belief that they give “reasonable” performance for many problems.

We did not know in advance how the generators would perform with the selected parameters on any particular problem in the experiments. The parameters were selected without regard for the particular problems included in the experiments.

In addition to the UNIX command-line arguments for the search limits given above, the following arguments were also used to invoke the generators.

`-postkind pi -negtkind pi -randseed -one -bpnone -bw`

The `postkind pi` and `negtkind pi` arguments specify that Π and the tuple shifts are as described earlier. The `randseed` argument means that each time the generator is invoked, it uses a new seed for producing pseudorandom numbers. Thus, on successive executions, the generator may follow different execution paths. Executing the generator many times gives a sample of the generator’s possible behavior over the possible random number sequences. The `bw` argument indicates the use of black and white tuples (as opposed to all-black or all-white tuples). The `one` and `bpnone` arguments are instructions related to backtracking; their precise meaning is unimportant here.

3 Description of Experiment Data

The data is presented in Appendix A. There is one table for each (dichotomy, n , generator) triple. All times are in CPU seconds.

It is convenient to explain the data by example, so we refer to Table 1. At the top of the table is the mean execution time for the $p = 1, q = n - 1$ case; here, .2132 CPU seconds. The $p = 1, q = n - 1$ case is typically the easiest (p, q) pair to find tuples for. The times within the table are scaled to this case, to clarify the changing difficulty of finding tuples as we move away from $p = 1, q = n - 1$.

Within the table, the nonblank table entries can be categorized into three types (not including shading). The first type of entry has two numbers stacked above each other, as with $p = 1, q = 2$. The $p = 1, q = 2$ entry indicates that no failures occurred for $p = 1, q = 2$; the average scaled time of the 50 successes was 103 (103 times the base time of .2132 seconds), and the standard deviation of the scaled time was 73. Another type of table entry has five numbers. For example, with $p = 10, q = 3$, the mean scaled time for the 50 successes was 544, with a standard deviation of 395. There were 42 failures; the average scaled time for these 42 failures was 1313; and the standard deviation was 37. A third type of table entry occurs when there are no successes, as with $p = 2, q = 2$. For $p = 2, q = 2$ there were 51 failures; the average failure time was 1074, with a standard deviation of 24.

When available, the tables show the optimal (largest) p values for a given value of q . Optimal values are indicated by dark shading in the appropriate table entry. (No trials were conducted for certain optimal (p, q) pairs; this happens when the generators did not achieve optimality, or when the optimal p value is larger than 10 and is not a multiple of 5.) When it is known that no solution exists for a given value of q , this is indicated by an asterisk next to the column header. The optimality of the displayed values was determined by exhaustive search. In some cases, when it was not possible to determine the optimal values, the best known solution values are indicated with light shading. For example, see Table 2.

Considerable CPU time was used in the attempt to determine optimal values. In some cases, weeks of CPU time were spent determining whether a solution exists for a given (p, q) pair.

The procedure for determining optimal values was independent of the C-like pseudocode in Section 2. This procedure followed the staircase pattern described in [1]. For example, in Table 1, to find the optimal values, the generator traversed (p, q) pairs in the following order: $(1, 2)$, $(2, 2)$, $(2, 3)$, $(3, 3)$, $(4, 3)$, $(5, 3)$, \dots , $(11, 3)$, $(12, 3)$. The optimal pairs are therefore $(1, 2)$ and $(11, 3)$, since no solutions were found for $(2, 2)$ or $(12, 3)$.

The tables were produced by a *perl* script that scans the 50 to 101 generator output files for each relevant (p, q) pair in a given table, and automatically produces the L^AT_EX source to make the table.

4 Discussion

In this section we look at the three questions that the experiments were designed to answer. In Section 4.4, we consider the quality of solution tuples as a function of the dichotomy.

4.1 Tradeoff Between Difficulty and Execution Time

Apparently, solution tuples are relatively abundant when p, q is close to $p = 1, q = n - 1$. As we move away from this point, ostensibly the problem becomes increasingly constrained and the solution tuple density decreases. For a given dichotomy and value of n , there is a tradeoff between execution time and the difficulty of the problem as defined by the (supposed) solution density. Inspection of the tables shows that Gen0 and Gen1 require more execution time as we move away from $p = 1, q = n - 1$. Thus, we can control the amount of time to find tuples by selecting p and q .

4.2 Time to Find Tuples

Here we address the issue of whether or not there exists an algorithm that finds distinguishing tuples reasonably quickly, when such tuples exist for a given dichotomy and value of n . Inspection of the tables in Appendix A yields an affirmative answer, at least for the problems in these experiments.

When solutions are relatively abundant or when solutions exist in a constrained situation such as for small n , either Gen0 or Gen1 is sufficient to find tuples. For the other cases, Gen0 does not do as well as Gen1.

4.3 Gen1 vs. Gen0

The third purpose of the experiments was to show that Gen1 is a “good” algorithm by comparing it to the benchmark algorithm Gen0.

For the most part, the measure of an algorithm should be how it does on difficult problems. (It is easy to do well on easy problems.) The definition of a difficult problem depends in general on the algorithm used to solve it. Here, we say that the difficulty increases when n increases and when (p, q) moves away from $p = 1, q = n - 1$. One reason for this is that (empirically) these cases take more time; in every table in Appendix A, as we move away from $p = 1, q = n - 1$ we see increased time to find a solution. Another reason is that the tuples away from $p = 1, q = n - 1$ are the ones that are most desirable for OCR [1].

Examination of the tables shows that for the difficult problems as defined above, Gen1 finds solutions faster and over a wider range than Gen0. For the “easy” problems, Gen0 does better.

It is instructive to examine why Gen0 does better than Gen1 on easy problems. One reason is that Gen1 was designed with difficult problems in mind. A consequence of this design is that Gen1 cannot find any solution without evaluating many tuples at each search node, even when the

search goes directly to a solution with no backtracking. On the other hand, Gen0 only evaluates one tuple at each search node. These facts and the experimental data suggest that we might want to form a hybrid algorithm that executes Gen0 and Gen1 in parallel. If t_0 and t_1 are the respective execution times for Gen0 and Gen1 on a given problem, then the hybrid algorithm takes time that is at worst roughly $2 \min(t_0, t_1)$. This is in some ways more desirable than the execution characteristics of either generator. For easy instances, the hybrid uses at most twice the time of Gen0; for difficult instances, the hybrid uses at most twice the time of Gen1.

There are easy instances in these experiments where Gen0 finds tuples over a wider range than Gen1 does, e.g., the *c-e* dichotomy for $n = 4$. For these instances, the failure times for Gen0 are on the order of 15 or 20 minutes, whereas the failure times for Gen1 are on the order of a few seconds. This suggests that for these instances the search width parameter of 20 used for Gen1 is too restrictive; this value causes Gen1’s search tree to be small and to contain no solutions. We suspect that Gen1 would find tuples for these instances if the search width were increased. The experiments appear, in these instances, to be unfair to Gen1, since Gen0 is allowed to spend more time searching.

4.4 Quality of Tuples as a Function of the Dichotomy

Here we briefly consider how, if at all, the quality of solution tuples depends on the dichotomy.

The best known values of p for a given value of q are summarized in Figure 2 ($n = 4$), Figure 3 ($n = 7$), and Figure 4 ($n = 10$), for each of the five dichotomies. The figures represent the shaded boxes that appear in the tables in Appendix A. Figure 2 ($n = 4$) represents optimal values; Figures 3 and 4 contain some entries that may not be optimal.

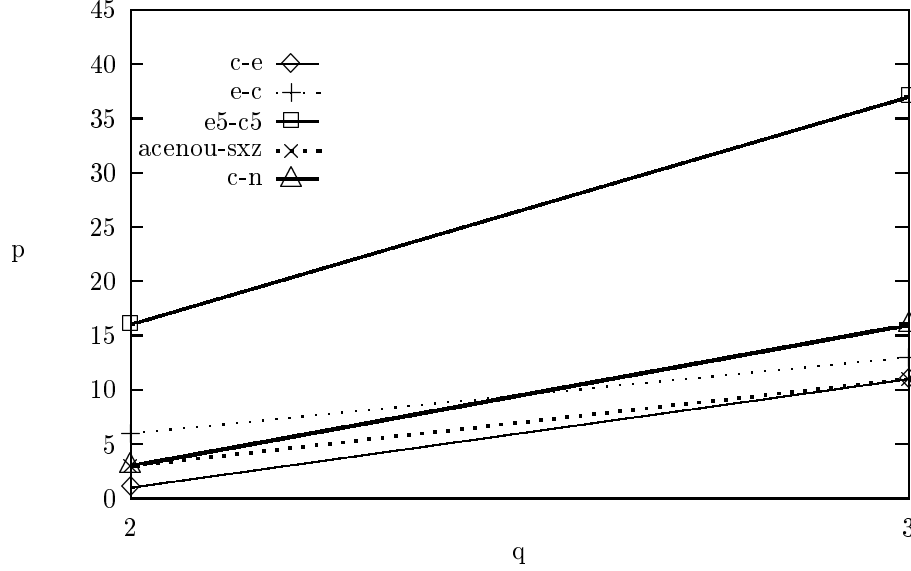


Figure 2: Optimal p for given q ($n = 4$).

A criterion for comparing the quality of tuples available to two dichotomies A and B is as follows: we can consider A to have higher quality tuples available than B if, in a plot of optimal values such as Figure 2, the curve for A is everywhere at or above the curve for B .

According to this criterion, the e_5-c_5 dichotomy has better tuples available than the other four dichotomies. Using Figures 2, 3, and 4, the five dichotomies can be roughly ordered in increasing order of quality of available tuples as follows:

$$c-e < acenou-sxz < e-c \leq c-n < e_5-c_5.$$

This ordering is consistent with our earlier belief that the $c-e$ dichotomy is a difficult one to find tuples for. The plots suggest that higher quality tuples are available for the $e-c$ dichotomy than for its converse (except when $q = n - 1$).

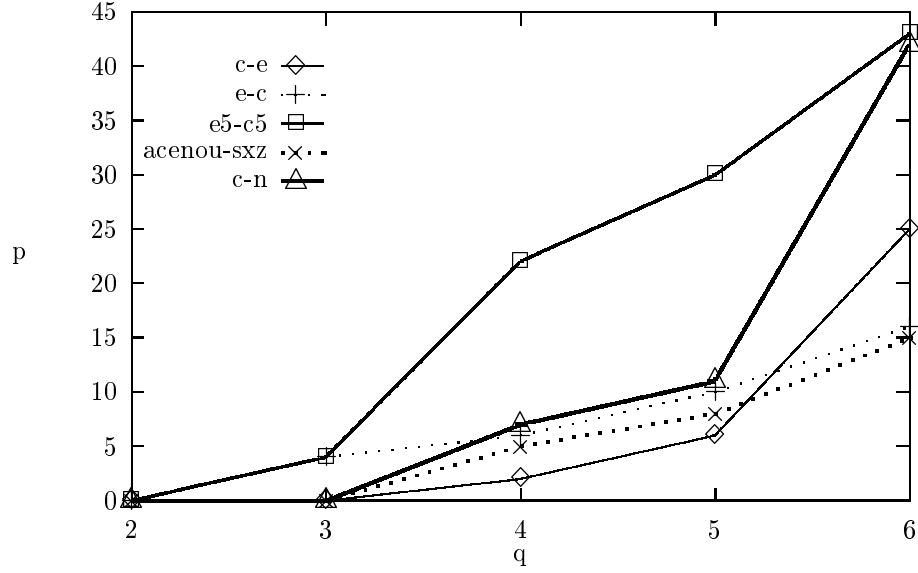


Figure 3: Largest known p for given q ($n = 7$).

The change in behavior at $q = n - 1$ is unexplained. It may be an artifact of the fact that the plots use only the best known values of p , and not necessarily the optimal values. This is another reminder that the data reported in this section is not exact and should be treated cautiously.

Despite their non-exact nature, the plots make it fairly clear that the quality of available tuples varies widely across dichotomies.

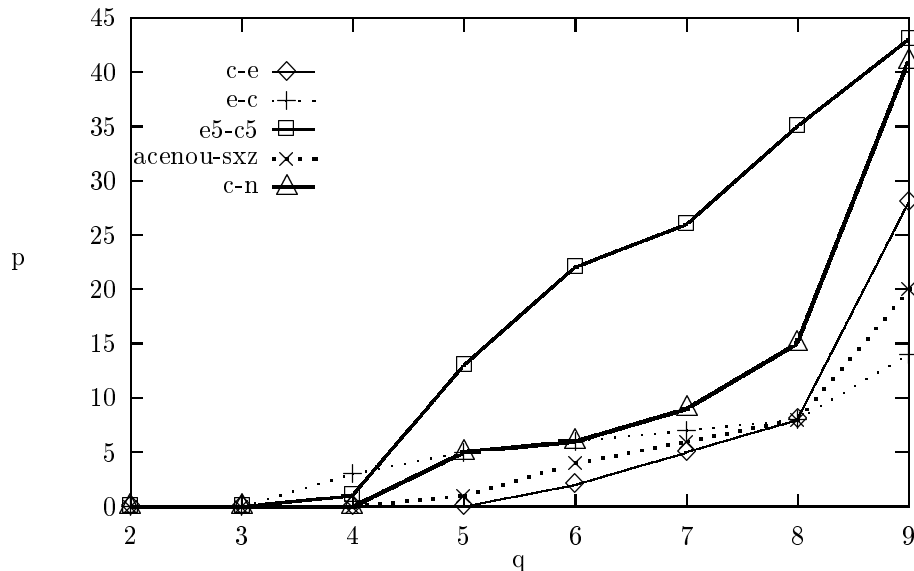


Figure 4: Largest known p for given q ($n = 10$).

5 Acknowledgement

I would like to thank D. Jung, M. Krishnamoorthy, and G. Nagy for commenting on a preliminary version of this report. G. Nagy participated in the design of the experiments, and suggested the plots in Section 4.4. Scanned characters were supplied by D. Jung.

References

- [1] D. Jung and M. Krishnamoorthy and G. Nagy and A. Shapira, “N-Tuple Features for OCR Revisited,” Submitted to IEEE Transactions on Pattern Analysis and Machine Intelligence, 1995.

Appendix A: Experiment Data

This appendix contains 30 tables, one for each (dichotomy, n , generator) triple.

Note: to reduce the table width, Table 21 does not have a column for $q = 3$. The data that would go in this column is that for $p = 1, q = 3$ there were 51 failures with an average scaled time of 140 and a standard deviation of 2.5. Exhaustive search determined that there are no solutions for $q = 3$.

Mean time for p=1,q=3: .2132 s

	<i>q</i>	
	2	3
1	103 73	1 0.08
2	1074 ⁵¹ 24	1.0 0.17
3		1.6 1.2
4		2.6 3.2
5		4.1 3.8
6		7.6 7.8
7		23 22
8		149 230
9		182 255
<i>p</i>		
10		544 1313 ⁴² 395 37
11		
15		1296 ⁵¹ 40
20		
25		
30		
35		
40		

Table 1: Success and failure times, n = 4, c-e, Gen0.

Mean time for p=1,q=6: .169 s

	q					
	2*	3*	4		5	6
1		1390 ⁵¹ 89	327 345	1328 ³³ 80	2.7 3.4	1 0.19
2			1338 0	1330 ⁵¹ 132	57 119	1.1 0.26
3					185 246	1495 ⁵ 206
4					375 382	1366 ²⁶ 154
5					467 482	1403 ⁵¹ 174
6						43 75
7						1083 ⁴ 87
8						100 223
9						1153 ¹¹ 143
p						172 276
10						1230 ⁹ 247
15						186 281
20						1211 ³⁰ 251
22						298 281
25						1314 ⁴⁸ 229
30						273 152
35						1246 ⁵¹ 198
40						

Table 2: Success and failure times, $n = 7$, c-e, Gen0.

Mean time for p=1,q=9: .2066 s

	q							
	4	5	6	7	8	9		
1			442 1283 ⁵¹ 352 186	113 1354 ³ 261 106	2.2 4.4	1 0.097		
2				307 1347 ⁴⁸ 336 293	16 1049 ¹ 51 0	1.0 0.09		
3				448 1345 ⁵¹ 364 281	130 1509 ²⁶ 214 334	3.7 1652 ⁴ 15 107		
4					260 1296 ⁵¹ 343 306	43 1283 ⁵ 263 325		
5						68 1838 ⁹ 229 149		
6						125 1396 ¹⁶ 321 331		
7						125 1388 ²⁸ 312 324		
8						163 1202 ³² 318 302		
9						107 1375 ⁵¹ 190 350		
p								
10								
15								
20								
25								
28								
30								
35								
40								

Table 3: Success and failure times, $n = 10$, c-e, Gen0.

Mean time for $p=1, q=3$: .1712 s

	q	
	2	3
1	6.9 6.9	1 0.13
2	8.1 6.3	0.99 0.13
3	8.3 5.8	1.0 0.16
4	6.7 5.5	0.98 0.17
5	7.6 5.2	1.0 0.16
6	37 39	1.1 0.16
7	1190 ⁵¹ 38	1.1 0.22
8		1.5 1.0
9		1.9 1.4
p		
10		4.7 4.2
13		
15		1362 ⁵¹ 52
20		
25		
30		
35		
40		

Table 4: Success and failure times, $n = 4$, e-c, Gen0.

Mean time for p=1,q=6: .2062 s

	q					
	2*	3		4	5	6
1	1277 ⁵¹ 19	174 238	1278 ¹⁰ 22	2.6 2.2	1.0 0.13	1 0.033
2		186 201	1228 ⁴² 30	18 28	1.9 2.8	4.3 23
3		413 306	1176 ⁵¹ 73	60 108	1215 ¹ 0	3.4 11
4				192 294	1302 ⁴ 18	12 25
5				333 385	1242 ²³ 123	23 48
6				418 411	1249 ⁵¹ 132	175 266
7					1057 ⁶ 119	27 123
8					378 362	1273 ⁴¹ 184
9					723 396	1313 ⁵¹ 190
p						42 93
10						1070 ⁶ 218
15						145 279
20						1173 ¹⁴ 243
25						212 329
30						1180 ²⁰ 285
35						1205 ⁵¹ 277
40						

Table 5: Success and failure times, $n = 7$, e-c, Gen0.

Mean time for p=1,q=9: .2076 s

	q											
	4		5		6		7		8		9	
1	290 364	1303 ⁵¹ 74	184 225	1242 ⁶ 122	11 24		13 83		1.0 0.16		1 0.04	
2			397 349	1182 ⁵¹ 95	109 225	1250 ⁵ 205	23 123	1078 ² 180	20 98		1.0 0.066	
3					248 223	1232 ²⁰ 177	79 175	1145 ⁷ 203	22 82	1548 ⁹ 246	3.9 15	
4					415 374	1244 ⁵¹ 215	217 309	1378 ²¹ 260	35 94	1328 ¹⁴ 268	11 33	1070 ¹ 0
5							306 274	1254 ⁵¹ 259	102 207	1207 ⁴⁰ 300	14 44	1435 ⁷ 345
6									191 230	1276 ⁵¹ 276	82 209	1231 ²² 380
7											144 263	1232 ³⁶ 379
8											198 267	1156 ⁵¹ 310
9												
p												
10												
14												
15												
20												
25												
30												
35												
40												

Table 6: Success and failure times, n = 10, e-c, Gen0.

Mean time for $p=1, q=3$: .2198 s

	q	
	2	3
1	1.1 0.13	1 0.044
2	1.2 0.23	0.98 0.083
3	1.3 0.29	1.0 0.054
4	1.5 0.48	1.0 0.06
5	1.8 0.65	1.0 0.057
6	2.1 1.3	1.0 0.073
7	2.1 1.1	0.99 0.083
8	2.2 1.3	1.0 0.12
9	3.1 1.8	0.99 0.071
p		
10	3.2 1.9	1.0 0.08
15	9.8 15	1.3 0.8
16		
20	1140 ⁵¹ 8.9	1.2 0.6
25		1.8 1.3
30		4.2 3.1
35		371 1263 ¹⁸ 391 30
37		
40		1223 ⁵¹ 48

Table 7: Success and failure times, $n = 4$, e_5 - c_5 , Gen0.

Mean time for p=1,q=6: .1924 s

	q				
	2*	3	4	5	6
1	1299 ⁵¹ 47	99 106	2.0 1.6	0.99 0.19	1 0.19
2		307 1156 ⁷ 300 76	4.9 11	1.0 0.2	0.99 0.18
3		511 1119 ⁵¹ 335 58	16 32	1.1 0.34	1.0 0.32
4			9.6 16	1.5 2.7	1.0 0.24
5			14 25	25 164	1.0 0.36
6			15 25	1.2 0.58	2.6 11
7			22 38	2.3 4.1	3.0 1268 ¹ 9.3 0
8			34 64	35 124	5.3 24
9			59 83	22 74	1.6 1320 ¹ 2.9 0
10			59 1071 ² 109 93	3.7 910 ¹ 5.4 0	5.3 1118 ² 18 179
15			278 1075 ¹⁰ 285 150	55 913 ¹ 171 0	65 1008 ⁴ 154 93
20			496 1047 ⁵⁰ 328 141	100 1021 ⁹ 166 220	53 1038 ¹⁰ 136 178
22					
25			1070 ⁵¹ 154	354 1086 ⁵¹ 254 189	83 1072 ¹³ 202 194
30					161 1067 ⁵¹ 213 220
35					
40					
43					
45					

Table 8: Success and failure times, $n = 7$, e_5 - c_5 , Gen0.

Mean time for p=1,q=9: .2212 s

	q											
	4		5		6		7		8		9	
1	500 254	1333 ⁵¹ 75	214 326	1313 ¹² 123	28 131		24 139		1.0 0.063		1 0.065	
2			314 362	1291 ²³ 118	46 146	1447 ¹ 0	22 116		2.9 10		1.0 0.086	
3			410 342	1197 ³⁸ 133	63 159	1197 ² 157	8.4 36	1527 ² 191	23 120		1.1 0.6	
4			559 448	1170 ⁵¹ 101	155 280	1328 ³ 95	15 65	1719 ² 63	13 55	1198 ⁴ 389	1.1 0.47	
5					68 147	1151 ⁹ 199	26 98	1491 ¹ 0	16 91	1485 ⁴ 245	1.3 1.2	1027 ⁴ 190
6					132 208	1136 ⁸ 236	28 68	1225 ⁵ 318	40 151	1194 ³ 201	1.7 4.0	1267 ³ 523
7					198 254	1119 ¹⁹ 156	33 65	1256 ⁷ 314	48 150	1287 ¹⁰ 253	5.6 32	1343 ³ 285
8					222 240	1102 ¹² 174	36 76	1123 ¹¹ 203	45 147	1080 ⁷ 242	2.3 5.2	999 ⁶ 235
9					168 190	1050 ⁴² 159	105 271	1115 ¹² 261	53 116	1168 ⁸ 226	8.4 30	1264 ⁶ 449
10					279 269	1091 ³¹ 174	67 128	1082 ⁷ 256	50 138	1185 ¹⁷ 290	42 140	1376 ⁹ 348
13												
p												
15					181 195	990 ⁵¹ 150	112 158	971 ⁵¹ 164	70 166	1244 ²¹ 343	18 74	1061 ¹⁸ 301
20									131 245	1154 ⁵¹ 302	78 193	1095 ⁴⁷ 277
22												
25											59 121	1014 ⁵¹ 232
26												
30												
35												
40												
43												
45												

Table 9: Success and failure times, $n = 10$, e_5 - c_5 , Gen0.

Mean time for $p=1, q=3$: .7872 s

	q	
	2	3
1	6.2 5.7	1 0.28
2	11 12	1.3 0.67
3	23 23	1.4 0.72
4	564 ⁵¹ 7.4	1.8 1.1
5		3.4 3.7
6		7.5 15
7		12 31
8		21 44
9		58 110
p		
10		137 835 ⁵ 185 26
11		
15		822 ⁵¹ 22
20		
25		
30		
35		
40		

Table 10: Success and failure times, $n = 4$, acenou-sxz, Gen0.

Mean time for p=1,q=6: 1.9652 s

	q							
	2*	3*	4		5		6	
1		324^{51} 42	136 121	345^{14} 27	9.7 23	272^2 14	1 3.0	
2			131 121	348^{21} 27	33 51	249^8 31	2.3 6.3	
3			138 114	338^{51} 34	64 80	279^{22} 70	6.3 15	236^4 57
4					89 98	282^{38} 68	18 46	236^9 53
5					114 115	313^{51} 69	25 48	276^{13} 53
6							35 53	252^{27} 47
7							51 72	263^{51} 50
8								
p								
9								
10								
15								
20								
25								
30								
35								
40								

Table 11: Success and failure times, n = 7, acenou-sxz, Gen0.

Mean time for $p=1, q=9$: 2.1424 s

	q					
	4	5	6	7	8	9
1				43 216 ⁵¹ 58 40	13 193 ¹⁴ 34 36	1 169 ¹ 2.6 0
2					39 217 ⁴¹ 52 42	6.6 219 ³ 18 68
3					57 203 ⁵¹ 67 45	14 186 ¹⁸ 33 50
4						18 180 ³⁸ 37 46
5						30 177 ⁵¹ 40 43
6						
7						
8						
p						
9						
10						
15						
20						
25						
30						
35						
40						

Table 12: Success and failure times, $n = 10$, acenou-sxz, Gen0.

Mean time for $p=1, q=3$: .1804 s

	q	
	2	3
1	12 8.6	1 0.17
2	31 35	0.98 0.17
3	63 62	1.0 0.2
4	1297 ⁵¹ 36	0.94 0.16
5		1.0 0.2
6		1.1 0.5
7		1.1 0.58
8		1.4 1.2
9		1.5 1.4
p		
10		2.4 2.2
15		154 221
16		
20		1292 ⁵¹ 90
25		
30		
35		
40		

Table 13: Success and failure times, $n = 4$, c-n, Gen0.

Mean time for p=1,q=6: .2064 s

	q				
	2*	3*	4	5	6
1		1199 ⁵¹ 65	21 28	1.2 0.35	1 0.077
2			93 131	1.6 1.5	1.0 0.12
3			140 1141 ²² 219 110	3.4 5.8	1.0 0.091
4			183 1200 ²⁰ 229 111	45 122	1.0 0.082
5			233 1168 ³¹ 281 109	58 1604 ¹ 132 0	2.1 6.1
6			402 1197 ⁵¹ 370 127	94 1101 ⁶ 169 148	1.3 1.0
7				114 1073 ⁸ 194 192	7.5 34
8				227 1127 ¹⁷ 303 234	8.2 30
9				430 1241 ⁵¹ 338 196	22 1114 ² 54 121
10					79 1368 ⁵ 270 284
p 11					
15					183 1372 ¹⁴ 234 315
20					332 1208 ⁵¹ 325 287
25					
30					
35					
40					
42					
45					

Table 14: Success and failure times, $n = 7$, c-n, Gen0.

Mean time for p=1,q=9: .1706 s

	q											
	4		5		6		7		8		9	
1			510 422	1242 ⁵¹ 136	125 239	1066 ⁶ 131	7.2 17	953 ¹ 0	1.0 0.28		1 0.16	
2					170 242	1154 ²⁶ 158	16 61	931 ¹ 0	1.3 1.2		0.95 0.13	
3					268 295	1147 ⁴⁶ 143	106 232	1123 ¹¹ 281	24 102	859 ¹ 0	1.0 0.42	
4					309 315	1186 ⁵¹ 178	161 255	1231 ²⁴ 253	52 113	1037 ³ 81	1.1 0.4	
5							155 251	1099 ⁵¹ 179	36 85	1049 ¹¹ 154	1.7 3.7	1109 ¹ 0
6									149 249	1177 ¹³ 282	4.1 17	954 ⁴ 101
7									184 302	1134 ³² 254	26 117	1119 ² 233
8									128 187	1215 ⁵¹ 341	34 118	1571 ⁴ 487
9											24 72	1039 ⁷ 182
p												
10											36 82	1199 ²³ 293
15											142 247	1272 ²⁹ 347
20											96 168	1297 ⁵¹ 363
25												
30												
35												
40												
41												
45												

Table 15: Success and failure times, n = 10, c-n, Gen0.

Mean time for p=1,q=3: .6368 s

	<i>q</i>	
	2	3
1	2.4 ⁵¹ 0.044	1 0.049
2		0.95 0.055
3		0.98 0.11
4		1.3 0.43
5		4.8 ⁵¹ 0.18
6		
7		
8		
9		
<i>p</i>		
10		
11		
15		
20		
25		
30		
35		
40		

Table 16: Success and failure times, n = 4, c-e, Gen1.

Mean time for p=1,q=6: .5668 s

	q				
	2*	3*	4	5	6
1		15 ⁵¹ 1.8	86 64	1.1 0.2	1 0.096
2			166 ⁵¹ 7.9	7.8 3.8	0.99 0.12
3				224 352 ⁵¹ 6.3 16	2.5 1.8
4					2.3 1.4
5					1.5 0.73
6					2.3 0.62
7					3.1 1.3
8					3.5 1.9
9					3.8 2.5
p					
10					5.4 4.1
15					94 467 ⁷ 66 27
20					457 ⁵¹ 17
22					
25					
30					
35					
40					

Table 17: Success and failure times, $n = 7$, c-e, Gen1.

Mean time for p=1,q=9: .8808 s

	q					
	4	5	6	7	8	9
1			83 191 ⁵¹ 3.0 7.1	0.98 0.04	1.0 0.041	1 0.054
2				33 221 ⁹ 23 16	5.4 2.6	0.92 0.078
3				171 ⁵¹ 13	5.3 0.59	2.2 1.4
4					6.9 3.0	2.2 1.4
5					8.1 158 ⁸ 9.0 9.6	1.1 0.47
6					150 ⁵¹ 8.9	1.9 0.42
7						2.4 0.97
8						3.0 1.3
9						2.2 1.1
p						
10						2.9 1.4
15						3.9 3.3
20						5.5 125 ² 4.0 3.5
25						109 ⁵¹ 10
28						
30						
35						
40						

Table 18: Success and failure times, $n = 10$, c-e, Gen1.

Mean time for $p=1, q=3$: .472 s

	q	
	2	3
1	1.2 0.14	1 0.12
2	1.2 0.15	1.1 0.12
3	1.2 0.14	1.0 0.14
4	1.2 0.12	1.0 0.11
5	3.0 ⁵¹ 0.29	1.0 0.1
6		1.1 0.13
7		1.2 0.12
8		1.2 0.13
9		1.2 0.18
p		
10		1.2 0.17
13		
15		4.6 ⁵¹ 0.44
20		
25		
30		
35		
40		

Table 19: Success and failure times, $n = 4$, e-c, Gen1.

Mean time for p=1,q=6: .6548 s

	q				
	2*	3	4	5	6
1	4.8 ⁵¹ 0.15	1.4 0.11	1.6 0.21	1.7 0.25	1 0.2
2		1.5 0.18	1.3 0.18	1.5 0.25	1.2 0.19
3		44 7.2	1.5 0.29	1.6 0.27	1.0 0.18
4		44 ⁵¹ 1.3	1.3 0.22	1.4 0.27	1.0 0.16
5			3.0 3.3	1.4 0.26	0.97 0.13
6			266 44	2.6 1.7	1.2 0.15
7			400 ⁵¹ 2.8	6.0 3.3	1.2 0.021
8				9.5 7.3	1.2 0.03
9				57 336 ⁵¹ 5.7 7.7	1.1 0.079
p					
10					1.0 0.029
15					203 ⁵¹ 7.2
16					
20					
25					
30					
35					
40					

Table 20: Success and failure times, $n = 7$, e-c, Gen1.

Mean time for p=1,q=9: .7346 s

	<i>q</i>					
	4	5	6	7	8	9
1	4.3 2.9	1.7 0.23	1.8 0.21	1.9 0.18	1.9 0.21	1 0.29
2	87 151 ⁴⁷ 43 7.1	1.4 0.22	1.4 0.18	1.6 0.24	1.5 0.21	1.2 0.27
3	267 ⁵¹ 9.8	1.9 0.72	1.7 0.23	1.7 0.24	1.6 0.26	0.96 0.23
4		42 334 ⁴ 67 9.2	1.4 0.21	1.4 0.2	1.4 0.18	1.1 0.21
5		378 ⁵¹ 19	2.9 3.3	3.1 3.6	1.3 0.18	0.98 0.17
6			415 ⁵¹ 20	50 416 ⁵¹ 18 19	2.0 1.7	1.4 0.036
7					4.9 3.6	1.3 0.017
8					8.1 3.3	1.2 0.079
9					318 ⁵¹ 12	1.1 0.13
<i>p</i> 10						1.0 0.052
14						
15						198 ⁵¹ 8.7
20						
25						
30						
35						
40						

Table 21: Success and failure times, n = 10, e-c, Gen1 (see note at beginning of Appendix A).

Mean time for $p=1, q=3$: .6522 s

	q	
	2	3
1	1.2 0.087	1 0.1
2	1.1 0.093	1.0 0.12
3	1.1 0.081	1.0 0.095
4	1.0 0.053	1.1 0.1
5	1.1 0.045	0.91 0.037
6	1.0 0.056	1.2 0.02
7	1.1 0.11	1.2 0.024
8	1.1 0.073	1.2 0.028
9	1.2 0.14	1.2 0.021
p		
10	1.3 0.18	1.2 0.021
15	1.4 0.1	1.1 0.021
16		
20	3.1 ⁵¹ 0.038	1.0 0.023
25		1.2 0.18
30		1.3 0.2
35		3.3 ⁵¹ 0.097
37		
40		

Table 22: Success and failure times, $n = 4$, e_5 - c_5 , Gen1.

Mean time for p=1,q=6: .6554 s

	q				
	2*	3	4	5	6
1	2.8 ⁵¹ 0.51	8.9 6.1	2.8 0.95	1.0 0.25	1 0.25
2		8.5 3.6	2.0 1.1	2.3 1.1	0.98 0.27
3		54 20	1.5 0.7	2.0 0.89	0.99 0.25
4		156 ⁵¹ 11	2.2 0.88	1.3 0.26	1.2 0.34
5			3.6 2.5	1.6 0.33	0.78 0.14
6			1.4 0.76	1.3 0.23	1.3 0.23
7			3.3 1.2	1.2 0.21	1.2 0.21
8			2.5 2.0	1.2 0.21	1.2 0.22
9			2.4 1.9	1.1 0.21	1.2 0.21
10			3.6 2.5	1.1 0.15	1.2 0.21
p 15			5.3 4.4	1.9 1.2	1.1 0.19
20			5.4 1.2	2.3 1.5	0.99 0.17
22					
25			56 ⁵¹ 8.5	6.3 8.9	1.7 0.97
30				145 ⁵¹ 8.4	3.4 123 ¹ 5.0 0
35					150 ⁵¹ 8.4
40					
43					
45					

Table 23: Success and failure times, $n = 7$, e_5 - c_5 , Gen1.

Mean time for p=1,q=9: .8522 s

	q					
	4	5	6	7	8	9
1	30 132 ⁵¹ 18 2.4	2.6 0.99	2.6 0.93	2.7 0.85	1.1 0.27	1 0.24
2		2.0 1.2	2.4 1.2	2.3 1.2	2.2 1.2	1.0 0.27
3		3.7 2.8	1.7 0.75	2.0 0.94	2.0 0.99	0.94 0.19
4		6.8 3.0	2.2 0.87	2.3 0.71	1.5 0.092	1.4 0.3
5		11 5.3	3.5 2.2	3.2 1.7	1.9 0.2	0.77 0.022
6		10 4.1	1.7 0.83	1.7 0.85	1.5 0.021	1.5 0.049
7		13 7.5	3.8 1.3	3.7 1.2	1.4 0.062	1.4 0.064
8		28 31	3.5 2.4	2.7 1.7	1.4 0.064	1.4 0.05
9		63 239 ⁹ 49 13	3.2 2.5	2.8 1.9	1.3 0.03	1.4 0.041
10		42 262 ⁴⁴ 19 20	4.6 3.0	4.1 2.7	1.3 0.047	1.4 0.055
13 p						
15		265 ⁵¹ 11	13 5.6	3.5 2.6	1.5 0.84	1.1 0.055
20			74 216 ⁵¹ 65 12	11 8.2	2.8 1.4	0.96 0.033
22						
25				130 ⁵¹ 15	2.6 132 ³³ 1.1 16	1.7 0.81
26						
30					118 ⁵¹ 5.5	3.0 126 ⁷ 3.5 6.0
35						144 ⁵¹ 6.8
40						
43						
45						

Table 24: Success and failure times, $n = 10$, e_5 - c_5 , Gen1.

Mean time for $p=1, q=3$: 1.6832 s

	q	
	2	3
1	1.0 0.18	1 0.15
2	1.1 0.2	0.93 0.093
3	1.1 0.24	0.89 0.099
4	2.5 ⁵¹ 0.4	0.89 0.12
5		0.95 0.14
6		0.91 0.15
7		0.96 0.15
8		0.9 0.15
9		0.9 0.16
p		
10		0.96 0.18
11		
15		5.4 ⁵¹ 0.75
20		
25		
30		
35		
40		

Table 25: Success and failure times, $n = 4$, acenou-sxz, Gen1.

Mean time for p=1,q=6: 2.235 s

	q				
	2*	3*	4	5	6
1		64 ⁵¹ 3.5	1.2 0.32	1.0 0.012	1 0.032
2			2.5 1.5	1.4 0.8	0.95 0.019
3			6.5 5.2	0.99 0.26	0.91 0.011
4			13 5.4	1.4 0.59	0.88 0.0089
5			12 3.0	2.0 1.3	1.8 1.2
6			110 ⁵¹ 1.5	5.1 6.5	0.87 0.028
7				135 ⁵¹ 2.9	0.87 0.0054
8					0.83 0.01
p					
9					0.81 0.008
10					0.97 0.3
15					422 ⁵¹ 6.1
20					
25					
30					
35					
40					

Table 26: Success and failure times, $n = 7$, acenou-sxz, Gen1.

Mean time for p=1,q=9: 2.216 s

	q					
	4	5	6	7	8	9
1		196 ⁵¹ 9.5	70 32	3.0 2.1	1.0 0.16	1 0.17
2			47 182 ⁵ 39 2.9	3.0 3.6	1.3 0.77	0.88 0.13
3			87 173 ¹² 45 11	14 9.0	1.1 1.0	0.89 0.14
4			186 ⁵¹ 19	9.3 7.3	1.7 0.56	0.84 0.14
5				13 3.6	1.8 1.1	1.5 1.2
6				139 ⁵¹ 6.1	6.0 135 ² 7.9 4.0	0.82 0.14
7					113 ⁵¹ 12	0.84 0.13
8						0.79 0.13
p						
9						0.75 0.13
10						0.88 0.3
15						264 ⁵¹ 17
20						
25						
30						
35						
40						

Table 27: Success and failure times, $n = 10$, acenou-sxz, Gen1.

Mean time for $p=1, q=3$: .5378 s

	q	
	2	3
1	1.9 0.38	1 0.078
2	2.8 ⁵¹ 0.44	0.98 0.084
3		0.98 0.092
4		0.98 0.086
5		0.95 0.093
6		0.95 0.089
7		0.95 0.13
8		1.1 0.16
9		1.0 0.18
p		
10		1.0 0.18
15		0.99 0.21
16		
20		6.4 ⁵¹ 1.1
25		
30		
35		
40		

Table 28: Success and failure times, $n = 4$, c-n, Gen1.

Mean time for p=1,q=6: .8056 s

	q				
	2*	3*	4	5	6
1		75 ⁵¹ 3.0	1.8 1.4	1.0 0.024	1 0.042
2			3.9 3.2	1.0 0.019	0.99 0.04
3			14 4.1	0.95 0.016	0.95 0.035
4			23 12	4.4 4.3	0.95 0.044
5			47 55	13 6.4	0.97 0.024
6			351 ⁵¹ 2.7	15 6.7	0.93 0.034
7				11 5.4	0.93 0.034
8				16 4.7	0.91 0.02
9				16 330 ¹⁴ 6.3 4.4	0.92 0.018
10				312 ⁵¹ 6.5	0.89 0.021
p 11					
15					2.7 1.8
20					13 6.3
25					100 92
30					429 535 ¹³ 99 11
35					232 414 ⁵¹ 161 13
40					
42					
45					

Table 29: Success and failure times, n = 7, c-n, Gen1.

Mean time for p=1,q=9: .8262 s

	q					
	4	5	6	7	8	9
1	173 ⁵¹ 26	70 144 ²⁸ 49 11	1.9 1.4	1.1 0.63	0.97 0.17	1 0.065
2		182 ⁵¹ 24	10 3.7	1.1 0.54	1.0 0.18	0.94 0.046
3			16 3.7	5.5 4.3	0.99 0.17	1.0 0.14
4			62 266 ¹³ 67 14	8.1 5.0	1.2 0.48	1.1 0.12
5			104 328 ⁴ 110 13	16 6.6	13 8.4	1.0 0.19
6			335 ⁵¹ 13	19 4.3	17 7.0	0.9 0.16
7				357 ⁵¹ 10	14 6.2	0.92 0.17
8					17 6.4	0.88 0.16
9					18 365 ²² 7.2 13	0.94 0.17
p 10					342 ⁵¹ 16	0.91 0.16
15						2.2 1.6
20						23 14
25						91 494 ²⁶ 60 20
30						346 435 ⁵¹ 91 27
35						
40						
41						
45						

Table 30: Success and failure times, $n = 10$, c-n, Gen1.